

Lecture 8 - Oct. 1

Math Review, Lab1 Solution, Lab2

Algebraic Properties of Relational Ops
Lab2: Celebrity_0
Functional Properties

Announcements/Reminders

- **Lab2** due this Friday
- Guide for **Programming Test 1** released

Relational Operations: Overriding

$$\{(x,y) \mid P_1(x,y) \vee P_2(x,y)\} = \{(x,y) \mid P_1(x,y)\} \cup \{(x,y) \mid P_2(x,y)\}$$

$$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$$

Example: Calculate r overridden with $\{(a, 3), (c, 4)\}$

Hint: Decompose results to those in t's domain and those not in t's domain.

$$\begin{aligned} \text{relation } r \text{ is overridden by relation } t &= \{(d, r') \mid (d, r') \in t\} \cup \{(d, r') \mid (d, r') \in r \wedge d \notin \text{dom}(t)\} \\ &= \{(d, r') \mid (d, r') \in t\} \cup \{(b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\} \\ &= t \cup (\text{dom}(t) \triangleleft r) \end{aligned}$$

$$r \triangleleft \underbrace{\{(a, 3), (c, 4)\}}_t = \underbrace{\{(a, 3), (c, 4)\}}_t \cup \underbrace{(\{a, c\} \triangleleft r)}_{\text{dom}(t) \triangleleft r} = \boxed{\phantom{\{(a, 3), (c, 4)\}}}$$

$r = \{(\cancel{a}, 1), (b, 2), (\cancel{c}, 3), (\cancel{a}, 4), (b, 5), (\cancel{c}, 6), (d, 1), (e, 2), (f, 3)\}$

$(a, 3) \rightarrow (c, 4)$

Example: Calculate r overridden with $\{(a, 3), (c, 4)\}$

$b \in \text{Account}$
 $\leftrightarrow \mathbb{Z}$

basically r ,

except all pairs with first elements
in the domain of t should agree with t .

t

transfer.

$\text{changes} = \{(a1, b(a1) - \text{amt})$
 $(a2, b(a2) + \text{amt})\}$

THINK: transfer event from Lab1

Exercises: Algebraic Properties of Relational Operations

$$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$$

$$r[S]$$

$$S = \{a, b\}$$

Define the **image** of set s on r in terms of other relational operations.

Hint: What range of value should be included?

$$r[S] = \text{ran}(S \triangleleft r)$$

$$\subseteq \text{ran}(r)$$

goes a relation

✓ Define r **overridden with** set t in terms of other relational operations.

Hint: To be in t 's domain or not to be in t 's domain?

$$r \triangleleft t = t \cup (\text{dom}(t) \triangleleft r)$$

$$r \in \boxed{S} \leftrightarrow \boxed{T} \quad \{a, b, c, d\} \quad \{1, 2, 3, 4\}$$

$$s \subseteq S \quad t \subseteq T$$

Problem: reconstruct relation r
in terms of restriction & subtraction.

$$r \stackrel{?}{=} (s \triangleleft r) \cup (s \triangleleft r)$$

$$r \stackrel{?}{=} (r \triangleright t) \cup (r \triangleright t)$$

Examples

$$r = \{(a, 1), (b, 2), (c, 3)\}$$

$$s = \{b, d\}$$

$$t = \{1, 4\}$$

Lab2: Relational Operators

$$\tilde{K} = \{ (Mark, Alan), (Tom, Alan), (Tom, Mark) \}$$

$$id = \{ (Alan, Alan), (Mark, Mark), (Tom, Tom) \}$$

Among persons, there is a special one: the celebrity. We define the knows relation which captures how the celebrity and other persons should be treated differently: $c \in P$

1. No person (including the celebrity) knows himself or herself.
2. The celebrity knows no one.
3. Everyone know the celebrity.

CONTEXT Celebrity_c0

CONSTANTS

k knows relation

c celebrity

P Set person

AXIOMS

$$axm1: P \subseteq \mathbb{N}$$

$$axm2: c \in P$$

$$axm3: k \in (P \setminus \{c\}) \leftrightarrow P$$

$$axm4: k^{-1}[\{c\}] = P \setminus \{c\}$$

$$axm5: k \cap id = \emptyset$$

END

MACHINE Celebrity_0

SEES Celebrity_c0

VARIABLES

r

INVARIANTS

$$inv1: r \in P$$

EVENTS

Initialisation

begin

$$act1: r := c$$

end

Event celebrity (ordinary) $\hat{=}$

begin

$$act1: r := c$$

end

END

$$k \in P \leftrightarrow P$$

$$k \in \mathcal{P}(P \times P)$$

$$e.g. \underline{P} = \{ Alan, Mark, Tom \}$$

$$\underline{k} = \{ (Alan, Mark), (Alan, Tom), (Mark, Tom) \}$$

$$\subseteq P \times P$$

$c \notin \text{dom}(k)$

model the set of persons via unique ids.

identity relation

$$axm5: \forall (p1, p2). (p1, p2) \in k \Rightarrow p1 \neq p2$$

$$\text{Axiom 4: } \underbrace{K^{-1}[\{c\}]}_{\text{celebrity is known by}} = \underbrace{P \setminus \{c\}}_{\substack{\text{tom} \\ \text{every person} \\ \text{except the} \\ \text{celebrity themselves}}}$$

$$K^{-1} = \{ (Mark, Alan), (Tom, Alan), (Tom, Mark) \}$$

$$\boxed{K^{-1}[\{Tom\}]} = \boxed{\{Alan, Mark\}}$$

celebrity Tom is known by everyone except Tom

Axiom 3 $k \in (P \setminus \{c\}) \leftrightarrow P$

III

$$k \in P \leftrightarrow P$$

$$c \notin \text{dom}(k)$$

Functional Property

$$r \in S \leftrightarrow T$$

isFunction(r) \Leftrightarrow

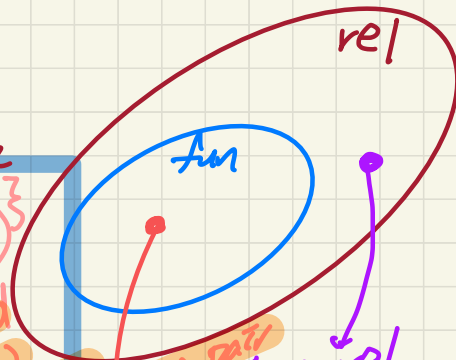
$\forall s, t1, t2 \bullet$

$$(s \in S \wedge t1 \in T \wedge t2 \in T)$$

\Rightarrow

$$((s, t1) \in r \wedge (s, t2) \in r \Rightarrow t1 = t2)$$

e.g. $\{ (a, 1), (b, 2), (a, 3) \}$



$(s, t1)$ and $(s, t2)$ refer to the same pair
two pairs sharing the same 1st value
2nd values are the same
a rel and a function at the same time.
some rel not qualified to be a fun.

the same dom value cannot map to distinct values in the range.

Q: Smallest relation satisfying the functional property. \emptyset

Q: How to **prove** or **disprove** that a relation r is a function.

Q: Rewrite the functional property using **contrapositive**.

$$p \Rightarrow q \equiv \neg q \Rightarrow \neg p$$

$$r = \{ (\underline{a}, \boxed{1}), (b, z), (\underline{c}, \boxed{1}) \}$$

1. Exercise: check this against the func. property
→ satisfied

2. injective function